## <u>Join index</u>: values of attribute index the fact table

<u>Hierarchically structured bit map index</u>

One level per attribute level

e.g. product  $\rightarrow$  categories



## **Example HOBI for products**



## Query 1

```
/* query Q1 */
SELECT COUNT(1)
FROM Auctions, Products, Categories, Days
WHERE Auctions.prodID = Products.prodID
AND Products.categID = Categories.categID
AND Auctions.dateID = Days.dateID
AND Categories.categName = 'Handheld'
AND Days.dateID
BETWEEN to_date('1-04-2007','dd-mm-yyyy');
AND to_date('31-07-2007','dd-mm-yyyy');
```

## Query 1

```
/* query Q2 */
SELECT Cities.cityName, SUM(Auctions.price)
FROM Auctions, Cities
WHERE Auctions.cityID = Cities.cityID
GROUP BY Cities.cityName;
```

## Query 1

/\* query Q3 \*/
SELECT Regions.regionName, SUM(Auctions.price)
FROM Auctions, Cities, Regions
WHERE Auctions.cityID = Cities.cityID
AND Cities.regionID = Regions.regionID
group by Regions.regionName;

## Time-HOBI

All dimensional models include Day, or a rollup of Day

Time-HOBI can eliminate a join to the Day dimension

Assumption: the fact table is sorted by date -reasonable and easily achieved It is typical for a fact table to stored in physical disk partitions based on date

Time-HOBI is HOBI augmented with a time index, TI

## Time Index (TI)

# Only one TI is created It will work with any/every HOBI index

The TI stores ranges of bit numbers belonging to a specified time interval

Each entry of a TI specifies a time interval and a range of bit numbers (bit positions in a bit map referencing the fact table)

## Time Index (TI)

# Implementation of TI (Section 5.3)

Tree structure where each level corresponds to a level of the day hierarchy:

```
Days \rightarrow Months \rightarrow Y ears
```

One tree per year - Each root is a year Next level is month Leaf level is day

## Time Index (TI)

# E.g. (based on figure 5)



Conceptual execution of query

Conceptually for any query that references time and attributes indexed by HOBI:

Using TI the system determines the fragments of the HOBI bit maps that need to be accessed.

Then, from the system retrieves HOBI fragments

The bit map fragments are 'combined' as required by the *where* clause

Lastly, the system retrieves rows of the fact table

# Query Q1: the number of auctions between April 1<sup>st</sup> and July 31<sup>st</sup>

SELECT count(1) from Auctions inner join Products ON ... inner join Categories ON ... inner join Days ON ... WHERE

*category name* = "Handheld"

AND

date between Apr 1, 2007 and July 31, 2007

To answer the above ... next slide

To answer the above:

- 1. the system determines the fragments to access that are relevant to "*between Apr 1, 2007 and July 31, 2007*"
- 2. the system retrieves the fragment for the Handheld category.
- 3. Because of "count(1)" there is no need to access Auctions ... the system can count the 1 bits

Q1-1 (related to Figure 5)

Query Q1 **Modified**: Consider Figure 5 and the number of auctions between April  $1^{st}$  and June  $30^{th}$  2009

SELECT count(1) from Auctions inner join Products ON ... inner join Categories ON ... inner join Days ON ... WHERE

```
category name = "Handheld"
```

AND

*date* between Apr 1, 2009 and Jun 30, 2009

- 1. To answer the above: the system determines the fragments to access that are relevant to "*between Apr 1, 2009 and June 30, 2009*" which are:
- 2. The system retrieves the fragment for the Handheld category: \_\_\_\_\_
- 3. The system counts the 1 bits: \_

Q1-2 (related to Figure 5)

Query Q1 **Modified**: Consider Figure 5 and the number of auctions in April and May of 2009

SELECT count(1) from Auctions inner join Products ON ... inner join Categories ON ... inner join Days ON ... WHERE

```
category name = "Handheld"
AND (
month = Apr
OR
month = May
AND
year = 2009
```

Processing required to answer the query .....